# Establishing Markov equivalence in cyclic directed graphs

Tom Claassen & Joris M. Mooij Radboud University Nijmegen, University of Amsterdam

> 39<sup>th</sup> Conference on UAI Pittsburgh, 2 August, 2023



Radboud University Nijmeger





- 2 The Cyclic Equivalence Theorem (CET)
- 3 An ancestral perspective on the CET
- 4 Establishing Markov equivalence and beyond

# Many important research questions are rooted in causality





#### benefits of exercise and healthy nutrition



racial and gender bias in AI



#### human activity and climate change



#### Covid vaccine efficacy

#### Ground truth causal model



## Meanwhile, in the real world ...



economic growth cycles





methane climate feedback

## The rabbit and the fox



rabbit :	$\frac{dx}{dt}$	=	$\alpha x - \beta x y$
fox:	$\frac{dy}{dt}$	=	$\delta xy - \gamma y$

Lotka-Volterra equations



periodic solution over time

## Static equilibrium solutions



rabbit: 
$$\frac{dx}{dt} = \alpha x - \beta x y$$
  
fox:  $\frac{dy}{dt} = \delta x y - \gamma y$ 

Lotka-Volterra equations



with `damping': equilibration towards unique, static solution





- **2** The Cyclic Equivalence Theorem (CET)
- 3 An ancestral perspective on the CET
- 4 Establishing Markov equivalence and beyond

## Linear/discrete cyclic causal models

$$X = \epsilon_X \quad (\sim \mathcal{N}(0, 1))$$
  

$$Y = \epsilon_Y \quad (\sim \mathcal{N}(0, 1))$$
  

$$Z = \alpha_{ZW}W + \alpha_{ZY}Y + \epsilon_Z \quad (\sim \mathcal{N}(0, 1))$$
  

$$W = \alpha_{WZ}Z + \alpha_{WX}X + \epsilon_W \quad (\sim \mathcal{N}(0, 1))$$



matching causal graph G

## **Key implication**

- directed global Markov property still holds [Spirtes, 1994; Bongers et al., 2021]
- standard *d*-separation still applies ...

linear cyclic Gaussian SCM

• ... but leads to a few extra quirks relative to acyclic models [Richardson, 1996]

#### Goal

• discovery aims for Markov equivalence class (cyclic PAG)



cyclic graph with two virtual edges

- A and D are *virtually adjacent* if they are not adjacent, but have a common child C in a cycle with A and/or D,
- an *itinerary* is a path over real and/or virtual edges (e.g. <*A*,*C*,*B*>)
- an itinerary is *uncovered* if no two nodes on the path are (virtually) adjacent, other than the neighbours along the path
- an itinerary <*A*,*C*,*D*> is a conductor if *C* is ancestor of *A* or *D*, otherwise it is a nonconductor,
- a nonconductor <*A*,*C*,*B*> is *perfect* if *C* is a descendant of a common child of *A* and *B*, otherwise it is *imperfect*

## Key cyclic terminology [Richardson, 1996/97]



cyclic graph where <A,D,F> and <D,F,B> are m.e. conductors w.r.t. uncovered itinerary <A,D,F,B>

 two triples <A,B,C> and <X,Y,Z> are mutually exclusive (m.e.) conductors w.r.t. an uncovered itinerary <A,B,C,..,X,Y,Z> if each consecutive triple along the itinerary is a conductor, all nodes are ancestor of each other but not of A or Z, and no two nodes are (virtually) adjacent, except along the itinerary itself Two directed graphs  $G_1$  and  $G_2$  are *d*-separation equivalent iff they have:

- i. the same (virtual) adjacencies,
- ii. the same unshielded conductors,
- iii. the same perfect nonconductors, (= v-structures')
- iv. the same m.e. conductors w.r.t. some uncovered itinerary,
- v. if  $\langle A, X, B \rangle$  and  $\langle A, Y, B \rangle$  are unshielded imperfect nonconductors in  $G_1$  and  $G_2$ , then X is ancestor of Y in  $G_1$  iff X is ancestor of Y in  $G_2$ ,
- vi. if <A,B,C> and <X,Y,Z> are m.e. conductors and <A,M,Z> is an unshielded imperfect nonconductor in G<sub>1</sub> and G<sub>2</sub>, then M is a descendant of B in G<sub>1</sub> iff M is a descendant of B in G<sub>2</sub>.



Example CET rule v : invariant edge D->E between two cycles C-D and E-F.

# Cyclic partial ancestral graphs (CPAGs)

• compact graphical representation to uniquely identify Markov equivalence class {*G*} for cyclic directed graph *G*, similar to standard (acyclic) PAGs

## **CPAG**

- edge between each pair of (virtually) adjacent nodes,
- arrowhead/tail marks to indicate invariant (non)ancestors, circle marks for noncommitted edge marks,
- **dashed-underlined**  $A \rightarrow \underline{B} \leftarrow C$  iff  $\langle A, B, C \rangle$  is an *imperfect nonconductor*.



Markov equivalent cyclic graphs

CPAG with dashed underlined triples





- 2 The Cyclic Equivalence Theorem (CET)
- **3** An ancestral perspective on the CET
- 4 Establishing Markov equivalence and beyond

#### Reflecting on the CET - CPAG

- the famous *Cyclic Causal Discovery* (CCD) algorithm [*Richardson,1996*] was an efficient CET-based implementation to reconstruct a CPAG from data,
- an analogous version based on *d*-separation could be used to establish Markov equivalence between cyclic graphs,
- yet, despite their central role in the CET, the dreaded `m.e. conductors on an uncovered itinerary' never need to be recorded explicitly in the CPAG ...

#### Motivating question

• does this mean we can simplify the CET? (spoiler: yes!)

#### Introducing the CMAG

**Key idea** (based on the familiar DAG-MAG-PAG trilogy from acyclic graphs):

- introduce the CMAG as intermediate ancestral representation,
- rephrase CET in terms of invariant elements in Markov equivalent CMAGs

The cyclic maximal ancestral graph (CMAG) *M* for directed graph *G* has:

- an edge between every pair of (virtually) adjacent nodes in G,
- a tail mark X —\* Y iff X is an ancestor of Y in G,
- arrowhead  $X \leftarrow * Y$  iff X is *not* an ancestor of Y,
- dashed-underlined  $A \rightarrow \underline{E} \leftarrow B$  for v-structures in M with a virtual edge in G.



directed cyclic graph

cyclic MAG

#### What about the m.e. conductors?

- In a CMAG *M*, quadruple  $\langle X, Z, Z', Y \rangle$  is a *u*-structure, iff there is an uncovered path  $X \rightarrow Z .. Z' \leftarrow Y$  in *M*.
- (Lemma 1) every *u*-structure in *M* matches a pair of m.e. conductors w.r.t. an uncovered itinerary in *G* and v.v.,
- comparing with virtual *v*-structures, both can be seen as arcs into a cycle

### Merge into one!

A triple  $\langle X, Z, Y \rangle$  is a virtual collider triple iff it is a virtual *v*-structure, or it is part of a *u*-structure  $\langle X, Z, Z', Y \rangle$  or  $\langle X, Z', Z, Y \rangle$ .



## Ancestral Cyclic Equivalence Theorem

Two CMAGs  $M_1$  and  $M_2$  are *d*-separation equivalent iff they have:

- i. the same skeleton,
- ii. the same *v*-structures,
- iii. the same virtual collider triples,
- iv. if  $\langle A, X, B \rangle$  and  $\langle A, Y, B \rangle$  are virtual collider triples in  $M_1$  and  $M_2$ , then X is ancestor of Y in  $M_1$  iff X is ancestor of Y in  $M_2$ .

 $\Rightarrow$  Basis for efficient procedure to establish Markov equivalence





- 2 The Cyclic Equivalence Theorem (CET)
- 3 An ancestral perspective on the CET
- 4 Establishing Markov equivalence and beyond

## Establishing Markov equivalence

Ancestral CET suggests straightforward Cyclic-Graph-to-CPAG procedure:

- convert directed graph G into CMAG M
- copy skeleton and *v*-structures in *M* to CPAG *P*
- use CMAG *M* to identify and orient virtual collider triples in CPAG *P*
- orient remaining edges between cycles via CET-rule (iv)
- compare resulting CPAGs

No need for *d*-separation tests!

- computational complexity scales as  $O(N^{2*}d^3)$
- worst case  $O(N^5)$  compared to previous  $O(N^7)$
- avg. scaling much better than worst case (for both versions)

# Example: Cyclic-Graph-to-CPAG procedure



(maximally informative) CPAG

## Experimental results – computational complexity



New ancestral perspective on the CET proved very useful:

- significantly simplified CET characterization
- helpful intermediate CMAG representation
- fast, graphical procedure to establish Markov equivalence

## Next steps

- recover maximally informative CPAG,
- merge *d*-separation (linear/discrete systems) and σ-separation (nonlinear systems) approaches via the CMAG,

## **Most promising**

 formulation via `virtual collider triples' suggests a natural extension of the CET to unobserved confounders, similar to the acyclic case, in the form of `virtual triples with order' [Ali et al,2009; Claassen&Bucur,2022]

# Thank you!

(poster 504 @11:00)