

A APPENDIX: PROOFS

Proof of Theorem 1. That $p'_{\epsilon,\ell}$ can be computed using the claimed number of evaluations of p follows immediately from Definition 5. Further, the approximation claim holds trivially for $r \leq g_\ell$ since p' and p are identical in this regime. In what follows, we argue that for any $r > g_\ell$, $p'_{\epsilon,\ell}(r)$, which is $p(g_j)$ for $j = \lfloor \log_{1+\epsilon} r \rfloor$, is a $(1 + \epsilon)$ -approximation of $p(r)$. By definition:

$$p(r) = \frac{p(g_j) g_j + \sum_{i=g_j+1}^r v(t_i)}{r}$$

Since each $v(t_i) \geq 0$, we have:

$$p(r) \geq \frac{p(g_j) g_j}{r} \geq \frac{p(g_j)}{1 + \epsilon}$$

where the second inequality can be verified by plugging in the definition of g_j for j as defined above. Thus, $p(g_j)$ is no more than $1 + \epsilon$ times $p(r)$.

For the other direction, we rely on the weak monotonicity assumption A.1. In particular, this implies $p(r) \leq p(g_j)$ whenever $r \geq g_j + m$, finishing the proof for such r . When $r < g_j + m$, however, $p(r)$ may be larger than $p(g_j)$. Nevertheless, we use $p(g_j + m) \leq p(g_j)$ and the definition of p to argue that $p(r)$ can be no more than $1 + \epsilon$ times $p(g_j)$ when $r < g_j + m$. To this end, consider the sequence of true labels $v(t_i) \in \{0, 1\}$ for $i = g_j + 1, g_j + 2, \dots, g_j + m$. Since $p(g_j + m) \leq p(g_j)$, it follows from the definition of q that at most a $p(g_j)$ fraction of these m labels is 1. Under these conditions, it can be verified that the largest possible value of $p(r)$ occurs at $r = r^* = g_j + \lfloor p(g_j) m \rfloor$ when all true labels $v(t_i)$ for $i = g_j + 1, \dots, r^*$ are 1. In this case:

$$\begin{aligned} \frac{p(r^*)}{p(g_j)} &= \frac{1}{p(g_j)} \frac{p(g_j) g_j + \lfloor p(g_j) m \rfloor}{g_j + \lfloor p(g_j) m \rfloor} \\ &\leq \frac{1}{p(g_j)} \frac{p(g_j) g_j + p(g_j) m}{g_j + p(g_j) m} = \frac{g_j + m}{g_j + p(g_j) m} \\ &\leq \frac{g_j + m}{g_j} = 1 + \frac{m}{g_j} \\ &\leq 1 + \frac{m}{g_\ell} = 1 + \frac{\lfloor \epsilon(1 + \epsilon)^\ell - 1 \rfloor}{\lceil (1 + \epsilon)^\ell \rceil} \\ &\leq 1 + \frac{\epsilon(1 + \epsilon)^\ell - 1}{(1 + \epsilon)^\ell} \\ &\leq 1 + \epsilon \end{aligned}$$

Thus, we have that for all $r < g_j + m$, $p(r)$ is at most $p(r^*)$ which is no more than $1 + \epsilon$ times $p(g_j)$. This finishes the proof. \blacksquare

Proof of Lemma 1. At each point g_k , we have s samples in the set X_k . Applying Eqn. (2), it follows that $q(g_k) =$

S_k/s , which is the average of the true labels of the s samples in X_k , provides a β -approximation for $p(g_k)$. \blacksquare

Proof of Theorem 2. By construction, $q(r) = p(r)$ in Algorithm 1 when $r \in \{1, \dots, \ell\}$. From Lemma 1, we further have $q(r)$ is a β -approximation of $p(r)$ for $r \in \{g_{\ell+1}, \dots, g_L\}$. It now follows from Corollary 1 that $step_{\epsilon,\ell}^q$ is a $\beta(1 + \epsilon)$ -approximation of p .

For the number of annotated samples used in the whole process, we have ℓ evaluations of v at points $1, \dots, g_\ell$ for computing $q(1), \dots, q(g_\ell)$, which subsumes the s samples needed for X_ℓ . For each of $X_{\ell+1}, \dots, X_L$ in subsequent steps, on average $\frac{\epsilon}{1+\epsilon}s$ new samples are drawn. This yields evaluations of v at a total of g_ℓ initial points and at $(L - \ell) \frac{\epsilon}{1+\epsilon}s = \frac{\epsilon(L - \ell)}{2(\beta - 1)^2(1 + \epsilon)p_{\min}^2} \ln \frac{L - \ell}{\delta/2}$ points chosen via stratified sampling. \blacksquare

Proof of Lemma 3. By the definition of $p(g_k)$, we have:

$$\begin{aligned} g_k p(g_k) &= \sum_{i=1}^{g_k} v(t_i) \\ &= \sum_{i=1}^{g_\ell} v(t_i) + \sum_{j=\ell}^{k-1} \sum_{i=g_j+1}^{g_{j+1}} v(t_i) \\ &= g_\ell p(g_\ell) + \sum_{j=\ell}^{k-1} \sum_{i=g_j+1}^{g_{j+1}} v(t_i) \end{aligned}$$

As observed earlier, $g_{j+1} - g_j \geq (1 + \epsilon)^{j+1} - (1 + \epsilon)^j - 1 = \epsilon(1 + \epsilon)^j - 1$. When $j \geq \ell$, this implies $g_{j+1} - g_j \geq \epsilon(1 + \epsilon)^\ell - 1 \geq m$. From the strong monotonicity assumption A.2, the inner summation may be bounded as:

$$p_\Delta(g_{j+1}) \leq \frac{\sum_{i=g_j+1}^{g_{j+1}} v(t_i)}{g_{j+1} - g_j} \leq p_\Delta(g_j)$$

Plugging these bounds back into the summation over j , we obtain:

$$\begin{aligned} g_k p(g_k) &\geq g_\ell p(g_\ell) + \sum_{j=\ell}^{k-1} (g_{j+1} - g_j) p_\Delta(g_{j+1}) \\ &= Y_-(\epsilon, \ell, k, \Delta) \\ g_k p(g_k) &\leq g_\ell p(g_\ell) + \sum_{j=\ell}^{k-1} (g_{j+1} - g_j) p_\Delta(g_j) \\ &= Y_+(\epsilon, \ell, k, \Delta) \end{aligned}$$

This finishes the proof. \blacksquare

Proof of Lemma 4. Using the same argument as in the proof of Lemma 3, for any $j \geq \ell$, we have $g_{j+1} - g_j \geq$

m. Further, by the strong monotonicity assumption A.2, we have $p_\Delta(g_{j+1}) \leq p_\Delta(g_j)$. However, in general, $p_\Delta(g_{j+1})$ may be arbitrarily smaller than $p_\Delta(g_j)$, making it difficult to lower bound the ratio of L and U using the ratio of these local precision terms.

Similar to Ermon et al. (2013), we instead collect all ‘‘coefficients’’ of $p_\Delta(g_j)$ for each $j \in \{\ell, \ell+1, \dots, k\}$ in the expressions for L and U , resp., and bound each pair of corresponding coefficients. The computations here are more intricate than in prior work because of the non-integrality of ϵ ; the previous derivation was for the special case where ϵ is effectively 1.

Using the assumption $p(g_\ell) \geq p_\Delta(g_\ell)$, it can be verified that for the claim it suffices to show two properties:

$$\gamma g_\ell \geq g_{\ell+1} \quad (7)$$

$$\gamma (g_{j+1} - g_j) \geq g_{j+2} - g_{j+1} \quad (8)$$

for all $j \in \{\ell+1, \dots, k-1\}$. For the first property:

$$\begin{aligned} \frac{g_{\ell+1}}{g_\ell} &= \frac{\lceil (1+\epsilon)^{\ell+1} \rceil}{\lceil (1+\epsilon)^\ell \rceil} \\ &\leq \frac{(1+\epsilon)^{\ell+1} + 1}{(1+\epsilon)^\ell - 1} \\ &= 1 + \epsilon + \frac{2+\epsilon}{(1+\epsilon)^\ell - 1} \\ &\leq 1 + \epsilon + \frac{2+\epsilon}{\epsilon(1+\epsilon)^\ell - 1} \\ &\leq 1 + \epsilon + \frac{2+\epsilon}{m} = \gamma \end{aligned}$$

where the second inequality follows from the precondition $\epsilon \in (0, 1]$. For the second property:

$$\begin{aligned} \frac{g_{j+2} - g_{j+1}}{g_{j+1} - g_j} &= \frac{\lceil (1+\epsilon)^{j+2} \rceil - \lceil (1+\epsilon)^{j+1} \rceil}{\lceil (1+\epsilon)^{j+1} \rceil - \lceil (1+\epsilon)^j \rceil} \\ &\leq \frac{(1+\epsilon)^{j+2} - (1+\epsilon)^{j+1} + 1}{(1+\epsilon)^{j+1} - (1+\epsilon)^j - 1} \\ &\leq \frac{(1+\epsilon)^{\ell+2} - (1+\epsilon)^{\ell+1} + 1}{(1+\epsilon)^{\ell+1} - (1+\epsilon)^\ell - 1} \\ &= 1 + \epsilon + \frac{2+\epsilon}{\epsilon(1+\epsilon)^\ell - 1} \\ &\leq 1 + \epsilon + \frac{2+\epsilon}{m} = \gamma \end{aligned}$$

The second inequality follows from the observation that the ratio under consideration here is a decreasing function of j . This finishes the proof. ■

Proof of Theorem 3. By construction, $q^-(r) = q^+(r) = p(r)$ in Algorithm 2 when $r \in \{1, \dots, \ell\}$. From Lemma 2, we further have $q^+(r)$ and $q^-(r)$ are γ -approximations of $p(r)$ from below and above, resp., for

$r \in \{g_{\ell+1}, \dots, g_L$. It now follows from Corollary 1 that $step_{\epsilon, \ell}^{q^-}$ and $step_{\epsilon, \ell}^{q^+}$ are $\gamma(1+\epsilon)$ -approximations of p from below and above, resp.

That PAULA uses a total of $g_\ell + \Delta(L - \ell)$ evaluations of v follows from the observation that it computes p exactly at points $1, \dots, g_\ell$, which in total require g_ℓ evaluations, and p_Δ at $L - \ell$ additional points, each of which requires Δ evaluations. ■